

# **bitelist.sty**

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## “Splitting” a List at a List Inside in T<sub>E</sub>X’s Mouth\*

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### **Abstract**

**bitelist.sty** provides commands for “splitting” a token list at the first occurrence of a contained token list (i.e., for given  $\sigma$ ,  $\tau$ , return  $\beta$  and shortest  $\alpha$  s.t.  $\tau = \alpha\sigma\beta$ ). As opposed to other packages providing similar features, (i) the method uses T<sub>E</sub>X’s mechanism of reading delimited macro parameters; (ii) the splitting macros work by pure expansion, without assignments, provided the macro doing the search has been defined before processing (e.g., a file); (iii) instead of using one macro for a “substring” test and another one to replace the “substring”—which includes extracting corresponding prefix and suffix—the same macro that detects the occurrence returns the split; (iv)  $\varepsilon$ -T<sub>E</sub>X is not required. (And L<sub>A</sub>T<sub>E</sub>X is not required.)

This improves the author’s **fifinddo.sty** (v0.51—and may once be used there). An elaborated approach (additionally to a simpler one) is provided that does not loose outer braces of prefix/suffix.

“Substring” detection and “string” replacement are (implicitly) included with respect to certain representations of characters by tokens. Counting occurrences and “global” replacement could be achieved by applying the operation to earlier results, etc.—so this approach seems to be “fundamental” for a certain larger set of list analysis tasks.

The documentation aims to prove the correctness of the methods with mathematical rigour.

**Related packages:** `datatool`, `stringstrings`, `ted`, `texapi`, `xstring`

**Keywords:** macro programming, text filtering, substrings

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\*This document describes version [v0.1](#) of **bitelist.sty** as of 2012/03/29.

<sup>†</sup><http://contact-ednotes.sty.de.vu>

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## 1 Task, Background Reasoning, and Usage

### 1.1 The Task Quite Precisely

Perhaps I should not have written “splitting” before, see Section 7 why I did so though. Actually:

At first we are dealing with token lists  $\tau$  and  $\sigma$  without braces (unless their category code has been changed appropriately) that can be stored as macros without parameter or in token list registers. We want to find out whether  $\tau$  contains  $\sigma$  (“as a subword”) in the sense that there are such token lists  $\alpha$  and  $\beta$  that  $\tau$  is composed as  $\alpha\sigma\beta$ , i.e.,

$$\tau = \alpha\sigma\beta$$

and in this case we want to get  $\alpha$  and  $\beta$  of this kind with  $\alpha$  being the *shortest* possible. I.e., if there are such  $\gamma$  and  $\delta$  that  $\tau$  is composed as  $\gamma\sigma\delta$ ,  $\alpha$  must be

contained as a “prefix” in  $\gamma$ , i.e.,  $\gamma$  is composed as  $\alpha\eta$  for some token list  $\eta$ . The token lists  $\alpha, \beta, \gamma, \delta, \eta, \sigma$ , and  $\tau$  are allowed to be empty throughout.

The task will be extended for some braces in Section 4.

## 1.2 Idea of Solution

$\text{\TeX}$ ’s mechanism of expanding macros ( $\text{\TeXbook}$  Chapter 20) at least has a built-in mechanism to return such  $\alpha$  and  $\beta$  *provided*  $\tau$  contains  $\sigma$ . Define

```
\def<cmd>#1\sigma#2\theta{\langle replace-def\rangle}
```

where  $\theta$  must be a token list (maybe of a single token) that won’t occur in  $\tau$ .<sup>1</sup> This is a **limitation** of the approach: It works for sets of such  $\tau$  only that do not contain any of a small set of tokens or combinations of them. (**bitelist** will use  $\backslash\text{BiteSep}$ ,  $\backslash\text{BiteStop}$ , and  $\backslash\text{BiteCrit}$ , or any other three that can be chosen.)

On the other hand,  $\text{\TeX}$ ’s *category codes* ( $\text{\TeXbook}$  Chapter 7) can ensure this quite well. E.g., we may assume that input “letters” always have category code 11 (or 12, or one of them), and for  $\theta$  we can choose letters with *different* category codes such as 3. Without such tricks, you may often assume that nobody will input certain “silly” commands such as  $\backslash\text{BiteStop}$ . (But it may become difficult when you use a package for replacement macros for generating its own documentation …)

With a  $\langle cmd \rangle$  as defined above,  $\text{\TeX}$  will

expand  $\langle cmd \rangle \tau \theta$  to  $\langle replace \rangle$ ,

where  $\langle replace \rangle$  will be the result of replacing (a) all occurrences of #1 in  $\langle replace-def \rangle$  by  $\alpha$  as wanted and (b) all occurrences of #2 in  $\langle replace-def \rangle$  by  $\beta$  as wanted. I.e.,  $\langle cmd \rangle$  returns  $\alpha$  as its first argument and  $\beta$  as its second argument. The reason is that  $\langle cmd \rangle$ ’s first parameter is delimited by  $\sigma$  and the second one by  $\theta$  in the sense of The  $\text{\TeXbook}$  p. 203. Our requirement to get the *shortest*  $\alpha$  for the composition of  $\tau$  as  $\alpha\sigma\beta$  is met because  $\text{\TeX}$  indeed looks for the *first* occurrence of  $\sigma$  at the right of  $\langle cmd \rangle$ .

## 1.3 When We Don’t Know ...

When  $\sigma$  does *not* occur in  $\tau$  and we present  $\tau\theta$  to  $\langle cmd \rangle$  as before,  $\text{\TeX}$  will throw an error saying “Use of  $\langle cmd \rangle$  doesn’t match its definition.” When the purpose is “substring detection” only, without returning  $\beta$ , many packages have solved the problem by issuing something like

```
\langle cmd \rangle \tau \sigma \theta
```

---

<sup>1</sup>I am still following others in confusing source code and tokens. I have better ideas, but must expand on them elsewhere. Writing  $\backslash\text{def}$  rather indicates that it is source code, then  $\sigma$  etc. should be replaced by strings that are converted into tokens  $\sigma$  etc.  $\langle cmd \rangle$  sometimes is a *string* starting with an escape character, or it is an active character; but sometimes it rather is an “active” *token* converted from such an escape string or an active character.

Then (still provided  $\theta$  does not occur in  $\tau$ )  $\langle cmd \rangle$ 's second argument is empty *exactly* if  $\sigma$  occurs in  $\tau$ . This method has, e.g., been employed in L<sup>A</sup>T<sub>E</sub>X's internal `\in@` mechanism (e.g., for dealing with package options) and by the `substr` package. `datatool` has used the latter's substring test (for  $\sigma$ ) before calling a macro for replacing ( $\sigma$  by another token list, perhaps thinking of character tokens).

This way you get the wanted  $\alpha$  as the first macro argument immediately indeed. An obstacle for getting  $\beta$  is that  $\langle cmd \rangle$ 's *second* argument now contains an occurrence of  $\sigma$  that is not an occurrence in  $\tau$ . In `fifinddo.sty` I didn't have a better idea than using another macro to remove the “dummy text” from the second argument. I considered it an advantage as compared with `datatool` that *one* macro could do this for *all* replacement jobs, while `datatool` uses *two* macros with  $\sigma$  as a delimiter for each  $\sigma$  to be replaced.

But still, `fifinddo` has used *two* macros for each replacement, the extra one being for presenting  $\tau$  to  $\langle cmd \rangle$ , using a job identifier. This could be improved within `fifinddo`, but I could never afford to take the time for this.

## 1.4 The Trick

The solution presented here is not very ingenious, many students would have found it in an exercise for a math course. My personal approach was looking at `\GetFileInfo` from L<sup>A</sup>T<sub>E</sub>X's `doc` package. There they try to get *two* occurrences of a space token this way:<sup>2</sup>

```
\def\@tempb#1#2#3\relax#4\relax{%
```

and `\@tempb` is called as

```
\@tempb\relax?\relax\relax\relax
```

or with  $\tau = \langle list \rangle$

```
\@tempb\langle list \rangle\relax?\relax\relax\relax
```

The final `\relax` may not be removed, but for `doc` it doesn't harm. It harms for *me* when I don't want to have a `\relax` in a `.log` file list. `\empty` would be better, however ...

The idea is to use a *three*-parameter macro for that *single* occurrence of  $\sigma$ . We introduce a “dummy separator”  $\zeta$  (or  $\langle sep \rangle$ , `\BiteSep`) between  $\tau$  and the “dummy text” and a “criterion”  $\rho$  ( $= \langle crit \rangle$ , `\BiteCrit`) for determining occurrence of  $\sigma$  ( $= \langle find \rangle$ ) in  $\tau$  ( $= \langle list \rangle$ ). Neither  $\zeta$  nor  $\rho$  must occur in  $\tau$ . We will have definitions about as

```
\def\langle cmd \rangle#1\sigma#2\zeta#3\theta{\langle replace-def \rangle}
```

or

```
\def\langle cmd \rangle#1\langle find \rangle#2\langle sep \rangle#3\langle stop \rangle{\langle replace-def \rangle}
```

---

<sup>2</sup>We are undoubling the hash marks inside the definition text of `\GetFileInfo`.

and  $\tau$  will be presented with context

$\langle cmd \rangle \tau \zeta \sigma \rho \zeta \theta$  or  $\langle cmd \rangle \langle list \rangle \langle sep \rangle \langle find \rangle \langle crit \rangle \langle sep \rangle \langle stop \rangle$

This ensures that  $\langle cmd \rangle$  finds its parameter delimiters  $\sigma$ ,  $\zeta$ , and  $\theta$ , in this order.  $\sigma$  occurs in  $\tau$  exactly if the second argument of  $\langle cmd \rangle$  is  $\rho$ , and in this case the first occurrence of the second parameter delimiter  $\zeta$  delimits  $\tau$ . Then  $\langle cmd \rangle$ 's first argument is  $\alpha$ , and the second one is  $\beta$ , as wanted.

$\langle cmd \rangle$ 's third parameter is delimited by the final  $\theta$  (`\BiteStop`). When  $\sigma$  occurs in  $\tau$ ,  $\langle cmd \rangle$ 's third argument starts after the first of the two  $\zeta$ , so it is  $\sigma \rho \zeta$ . It is just ignored, this way  $\langle cmd \rangle$  removes all the “dummy” material after  $\tau$ . When  $\sigma$  does *not* occur in  $\tau$ , we ignore all of its arguments, and the macro that invoked  $\langle cmd \rangle$  must decide what to do next, e.g., keeping  $\tau$  elsewhere for presenting it to another parsing macro resembling  $\langle cmd \rangle$ .

## 1.5 Installing and Calling

The file `bitelist.sty` is provided ready, installation only requires putting it somewhere where TeX finds it (which may need updating the filename data base).<sup>3</sup>

Below the `\documentclass` line(s) and above `\begin{document}`, you load `bitelist.sty` (as usually) by

```
\usepackage{bitelist}
```

between the `\documentclass` line and `\begin{document}`; or by

```
\RequirePackage{bitelist}
```

within a package file, or above or without the `\documentclass` line. Moreover, the package should work *without* L<sup>A</sup>T<sub>E</sub>X and may be loaded by

```
\input bitelist.sty
```

Actually, using the package for macro programming requires understanding of pp. 20f. of The TeXbook. On the other hand, the package may be loaded (without the user noticing it) automatically by a different package that uses programming tools from the present package.

# 2 Implementation Part I

## 2.1 Package File Header (Legalize)

```

1  \def\filename{bitelist}           \def\filedate{2012/03/29}
2  \def\fileversion{v0.1} \def\fileinfo{split lists in TeX's mouth (UL)}
3  %% Copyright (C) 2012 Uwe Lueck,
4  %% http://www.contact-ednotes.sty.de.vu
5  %% -- author-maintained in the sense of LPPL below --

```

---

<sup>3</sup><http://www.tex.ac.uk/cgi-bin/texfaq2html?label=inst-wlcf>

```

6  %%
7  %% This file can be redistributed and/or modified under
8  %% the terms of the LaTeX Project Public License; either
9  %% version 1.3c of the License, or any later version.
10 %% The latest version of this license is in
11 %%     http://www.latex-project.org/lppl.txt
12 %% There is NO WARRANTY - this rather is somewhat experimental.
13 %%
14 %% Please report bugs, problems, and suggestions via
15 %%
16 %%     http://www.contact-ednotes.sty.de.vu
17 %%

```

## 2.2 Proceeding without L<sup>A</sup>T<sub>E</sub>X

Some tricks from Bernd Raichle's `ngerman.sty`—I need L<sup>A</sup>T<sub>E</sub>X's \Provides-Package for `fileinfo`, my package version tools. With `readprov.sty`, it issues `\endinput`, close conditional before:

```

18  \begingroup\expandafter\expandafter\expandafter\endgroup
19  \expandafter\ifx\csname ProvidesPackage\endcsname\relax \else
20      \edef\fileinfo{\noexpand\ProvidesPackage{\filename}%
21          [\filedate\space \fileversion\space \fileinfo]}
22      \expandafter\fileinfo
23  \fi
24  \chardef\atcode=\catcode`\@
25  \catcode`\@=11 % \makeatletter

```

Providing L<sup>A</sup>T<sub>E</sub>X's `\@firstoftwo` and `\@secondoftwo`:

```

26  \long\def\@firstoftwo #1#2{#1}
27  \long\def\@secondoftwo#1#2{#2}

```

## 2.3 Basic Parsing (No Braces)

`\BiteMake{\langle def\rangle}{\langle cmd\rangle}{\langle find\rangle}` provides the parameter text (TeXbook p. 203) for defining (by `\langle def\rangle`) a macro `\langle cmd\rangle` that will search for `\langle find\rangle`:

```

28  \def\BiteMake#1#2#3{\#1#2##1#3##2\BiteSep##3\BiteStop}

```

With `\BiteFindByIn{\langle find\rangle}{\langle cmd\rangle}{\langle list\rangle}`, you can use a `\langle cmd\rangle` (perhaps defined by `\BiteMake`) in order to search `\langle find\rangle` in `\langle list\rangle`. This is expandable as promised:

```

29  \def\BiteFindByIn#1#2#3{%
30      #2#3\BiteSep#1\BiteCrit\BiteSep\BiteStop}

```

Preparing a possible `\edef` as `\langle def\rangle`:

```

31  \let\BiteSep\relax \let\BiteStop\relax

```

And this is important in any case for correct testing of occurrence:<sup>4</sup>

```
32 \catcode`\Q=7 \let\BiteCrit=\catcode`\Q=11
```

Perhaps you could increase safety of tests by using something similar to the funny Q for `\BiteSep` and `\BiteStop`. However, this would additionally require reimplementations of the macros for keeping braces (Section 4) using `\edef`.

## 2.4 Simple Conditionals

By `\BiteMakeIfOnly{\langle def \rangle}{\langle cmd \rangle}{\langle find \rangle}`, you can make a command  $\langle cmd \rangle$  that with

```
\BiteFindByIn{\langle find \rangle}{\langle cmd \rangle}{\langle list \rangle}{\langle yes \rangle}{\langle no \rangle}
```

chooses  $\langle yes \rangle$  if  $\langle find \rangle$  occurs in  $\langle list \rangle$  and  $\langle no \rangle$  otherwise.

```
33 \def\BiteMakeIfOnly#1#2#3{\BiteMake{#1}{#2}{#3}{\BiteIfCrit{##2}}}
```

`\BiteIfCrit{\langle suffix \rangle}{\langle yes \rangle}{\langle no \rangle}` is the basic test for occurrence of  $\langle find \rangle$  in  $\langle list \rangle$ :

```
34 \def\BiteIfCrit#1{\ifx\BiteCrit#1\expandafter\@secondoftwo
```

If  $\langle cmd \rangle$ 's second argument—same as `\BiteIfCrit`'s first argument—is empty, `\BiteCrit` is compared with `\expandafter`, so  $\langle yes \rangle$  is chosen. That is correct, it happens when  $\langle find \rangle$  is a suffix of  $\langle list \rangle$ .

```
35 \else \expandafter\@firstoftwo \fi }
```

## 2.5 Passing Results Completely—No Braces

So the previous `\BiteMakeIfOnly` generates pure tests on occurrence, giving away information about prefix and suffix. It may be considered a didactical step fostering understanding of the following. When, by contrast

```
\BiteMakeIf{\langle def \rangle}{\langle cmd \rangle}{\langle find \rangle}
```

has been issued, a later

```
\BiteFindByIn{\langle find \rangle}{\langle cmd \rangle}{\langle list \rangle}{\langle list \rangle}{\langle yes \rangle}{\langle no \rangle} (*)
```

will expand to

```
\langle yes \rangle{\langle prefix \rangle}{\langle find \rangle}{\langle suffix \rangle}
```

if  $\langle list \rangle$  is composed as  $\langle prefix \rangle \langle find \rangle \langle suffix \rangle$  and  $\langle prefix \rangle$  is the shortest  $\alpha$  such that there is some  $\beta$  with  $\langle list \rangle = \alpha \langle find \rangle \beta$ . Otherwise, (\*) will expand to

```
\langle no \rangle{\langle list \rangle}
```

This gives all the information available. For actual applications, it may be too much, and the macro programmer may do something in between of `\BiteMakeIfOnly` and `\BiteMakeIf`:

---

<sup>4</sup>The idea for the “funny Q” is from the `ifmtarg` package.

```
36 \def\BiteMakeIf#1#2#3{%
37     \BiteMake{#1}{#2}{#3}##4##5##6{%
```

In the replacement text, we first do the same as with `\BiteMakeIfOnly`:

```
38 \BiteIfCrit{##2}{%
```

What follows is new.  $\langle cmd \rangle$ 's third argument is ignored. The fourth keeps the original  $\langle list \rangle$ .  $\langle yes \rangle$  is  $\langle cmd \rangle$ 's fifth and  $\langle no \rangle$  is its sixth argument.

```
39 {##5##1##3##2}%% if #3 in ##4
40 {##6##4}%% otherwise
41 }%
42 }
```

In (\*),  $\langle list \rangle$  has been doubled. That was no mistake. It is due to a shortcoming of `\BiteFindByIn`. With

```
\BiteFindByInIn{\langle find \rangle}{\langle cmd \rangle}{\langle list \rangle}{\langle yes \rangle}{\langle no \rangle}
```

you get the same result as with (\*):

```
43 \def\BiteFindByInIn#1#2#3{\BiteFindByIn{#1}{#2}{#3}{#3}}
```

**TODO** not sure about command names yet

## 3 Example Applications

### 3.1 Splitting at Space

This work actually arose from modifying `\GetFileInfo` as provided by L<sup>A</sup>T<sub>E</sub>X's `doc` package so that it would deal reasonably with “incomplete” file info—for the `nicefilelist` package. `\GetFileInfo` works best when the file info contains at least *two* blank spaces. But how many are there indeed?—And I wanted to do it *expandably*: while `\GetFileInfo` issues *definitions* of `\filedate`, `\fileversion`, and `\fileinfo`, date, version, and info should be passed as *macro arguments*.

`\BiteIfSpace` tries splitting at the next blank space passes results:

```
44 \BiteMake{\def{\BiteIfSpace}{ }##4##6{%
45     \BiteIfCrit{#2}{#5##1##2}{#6##4}}}
```

The difference to the `\BiteMakeIf` construction is that we do not pass  $\langle find \rangle$ , the space—it's not essential. (**TODO** names may change ...)

Now

```
\BiteFindByInIn{\BiteIfSpace}{\langle list \rangle}{\langle yes \rangle}{\langle no \rangle}
```

will pass prefix/suffix to  $\langle yes \rangle$  or  $\langle list \rangle$  to  $\langle no \rangle$ . If this is needed frequently, here is a shorthand `\BiteGetNextWord{\langle list \rangle}{\langle yes \rangle}{\langle no \rangle}`:

```
46 \def\BiteGetNextWord{\BiteFindByInIn{ }\BiteIfSpace}
```

See a test in `bitedemo.tex` (Section 6).

### 3.2 Splitting at Comma

... left as an exercise to the reader ...

## 4 Keeping Braces: Reasoning

Now we want to generalize task (Section 1.1) and solution (Section 1.4) for the case that  $\tau = \langle list \rangle$  has (balanced) braces (with category codes for argument delimiters), while  $\sigma = \langle find \rangle$  still has not (does not work with our method). So with  $\tau = \alpha\sigma\beta$ ,  $\alpha$  (“prefix”) or  $\beta$  (“suffix”) or both may contain braces. But we consider another restriction: braces must be balanced in  $\alpha$  and in  $\beta$ , we don’t try parsing inside braces (as opposed to the search for asterisks in Appendix D of The TeXbook).

According to TeXbook p. 204, when a macro  $\langle cmd \rangle$  finds an argument formed as  $\{\langle tokens \rangle\}$ , in  $\langle cmd \rangle$ ’s replacement text only  $\langle tokens \rangle$  is used, i.e., outer braces are removed. So when  $\alpha = \{\langle tokens \rangle\}$ , a parser  $\langle cmd \rangle$  as defined by our methods above will return  $\langle tokens \rangle$  instead of  $\{\langle tokens \rangle\}$ —likewise for  $\beta$ . We are now trying to *keep* outer braces in prefix/suffix by a more elaborate method.

The idea is to present  $\tau = \langle list \rangle$  with context<sup>5</sup>

```
 $\langle cmd \rangle \backslash empty \langle list \rangle \langle stop \rangle \langle sep \rangle \langle find \rangle \langle crit \rangle \langle sep \rangle \langle stop \rangle$ 
```

or in the notation of Section 1.4

```
 $\langle cmd \rangle \backslash empty \tau \theta \zeta \sigma \rho \zeta \theta$ 
```

Then, if  $\langle find \rangle$  occurs in  $\langle list \rangle$ , we must remove the  $\backslash empty$  from the prefix that we get with the earlier method (easy) and  $\langle stop \rangle$  from the suffix (tricky, similar problem recurs). Using old  $\theta$  for a new purpose works here because  $\langle cmd \rangle$  will look for  $\theta$  only when it has found  $\zeta$  before.

Mere testing for occurrence is not affected.

```
 $\backslash BiteMakeIfOnly \quad \text{and} \quad \backslash BiteFindByIn$ 
```

still can be used. We provide an improved version of

```
 $\backslash BiteMakeIf \quad (\backslash BiteMakeIfBraces)$ 
```

and of

```
 $\backslash BiteFindInIn \quad (\backslash BiteFindInBraces).$ 
```

---

<sup>5</sup>Perhaps I am confusing  $\backslash empty$  and the token list containing just  $\backslash empty$  here?

## 5 Implementation Part II

### 5.1 Keeping Braces

```
\BiteFindByInBraces{\langle find \rangle}{\langle cmd \rangle}{\langle list \rangle}{\langle yes \rangle}{\langle no \rangle}
```

varies \BiteFindByInIn according to the previous:

```
47 \def \BiteFindByInBraces#1#2#3{%
48   #2\empty#3\BiteStop\BiteSep#1\BiteCrit\BiteSep\BiteStop{#3}}
```

Such a  $\langle cmd \rangle$  can be made by \BiteMakeIfBraces{\langle def \rangle}{\langle cmd \rangle}{\langle find \rangle}:

```
49 \def \BiteMakeIfBraces#1#2#3{%
50   \BiteMake{#1}{#2}{#3}##4##5##6{%
51     \BiteIfCrit{##2}}%
```

$\langle no \rangle$  works as before. For  $\langle yes \rangle$ , first the \empty in the prefix is expanded for vanishing. \BiteTidyI and \BiteTidyII continue tidying.

```
52 {\expandafter \BiteTidyI %& if #3 in ##4
53   \expandafter{##1}}% %& prefix
```

Another \empty avoids that removal of \BiteStop in suffix by \BiteTidyII removes outer braces:

```
54 { \BiteTidyII \empty##2 }% %% suffix
55 { #3 }% %% find
56 { ##5 }% %% yes
57 { ##6{##4} }% %% otherwise
58 }%
59 }
```

\BiteTidyI{\langle prefix \rangle}{\langle suffix \rangle} first expands \BiteTidyII for removing \BiteStop in  $\langle suffix \rangle$ . \empty from \BiteFindByInBraces remains and is expanded next for vanishing. Finally, \BiteTidied reorders arguments for operation of  $\langle yes \rangle$ :

```
60 \def \BiteTidyI#1#2{%
61   \expandafter\expandafter\expandafter \BiteTidied
62   \expandafter\expandafter\expandafter {#2}{#1}}
63 \def \BiteTidyII#1\BiteStop{#1}
64 \def \BiteTidied#1#2#3#4{#4{#2}{#3}{#1}}
```

### 5.2 Leaving the Package File

```
65 \catcode`@=\atcode
66 \endinput
```

### 5.3 VERSION HISTORY

```

67  v0.1    2012/03/26  started
68      2012/03/27  continued, restructured
69      2012/03/28  continued, separate sections for "Mere Occurrence"
70          vs. ...; keeping braces, \BiteIfCrit
71      2012/03/29  proceeding without LaTeX corrected, restructured
72

```

## 6 Examples/Tests

You should find a separate file `bitedemo.tex` with examples. It may be run separately with `tex` (Plain TeX)—demonstrating that `bitelist` is “generic”, then finish by entering `\bye`. With “`latex_bitedemo.tex`”, end the job by entering `\stop`. **Expandability** is demonstrated by the `\BiteFind` commands running with `\typeout`.

---

```

\def\filename{bitedemo.tex} \def\filedate{2012/03/29}
\def\fileinfo{demonstrating/testing bitelist.sty (UL)}
\expandafter\ifx\csname ProvidesPackage\endcsname\relax \else
  \edef\bitedemolatexstart{%
    \noexpand\ProvidesFile{\filename}%
    [\filedate\space\fileinfo]%
    \noexpand\RequirePackage{bitelist}}
  \expandafter\bitedemolatexstart
\fi
\ifx\BiteMakeIf\undefined \input bitelist.sty \fi
\ifx\typeout\undefined
  \def\typeout{\immediate\write17}%
  \newlinechar'^^J
\fi
\def\splitted #1#2#3{>>#1|#2|#3<<}
\def\unsplitted#1{>>#1<<}
\def\spacetocomma#1#2{>>#1,#2<<}
\BiteMakeIfOnly {\def}{\occurseyesno}{no}
\BiteMakeIf {\def}{\noshowsplit}{no}
\BiteMakeIfBraces{\def}{\noShowSplit}{no}
\typeout{^^J
  \BiteFindByIn {no}{\occurseyesno}
  {bonobo}{YES!}{NO!}
  \BiteFindByIn {no}{\noshowsplit}
  {bonobo}{bonobo}{\splitted}{\unsplitted}
  \BiteFindByInIn{no}{\noshowsplit}
  {bonobo}{\splitted}{\unsplitted}
  ^^J
  \BiteFindByIn {no}{\occurseyesno}
}
```

```

{bobobo}{YES!}{NO!}
\BiteFindByIn {no}{\noshowsplit}
{bobobo}{bobobo}{\splitted}{\unsplitted}
\BiteFindByInIn{no}{\noshowsplit}
{bobobo}{\splitted}{\unsplitted}
~~J
\BiteFindByInBraces{no}{\noShowSplit}
{{bo}no{bo}}{\splitted}{\unsplitted}
~~J
\BiteGetNextWord{bo no bo}{\spacetocomma}{\unsplitted}
\BiteGetNextWord{bo nobo} {\spacetocomma}{\unsplitted}
\BiteGetNextWord{bonobo} {\spacetocomma}{\unsplitted}
~~J}
\endinput

```

---

## 7 The Package's Name

This package deals with  $\text{\TeX}$ 's expansion mechanism. In Knuth's metaphor, this is  $\text{\TeX}$ 's mouth. I am not entirely sure, I have never understood it, or I have understood it only for a few days or hours. However, the package deals with “Lists in  $\text{\TeX}$ 's Mouth” as described in Alan Jeffrey's 1990 TUGboat paper (Volume 11, No. 2, pp. 237–245).<sup>6</sup>

“Splitting” in title and abstract is an attempt to describe the package briefly without speaking Mathematicalese. It roughly refers to certain string functions in various programming languages<sup>7</sup> with “split” in their name. However, there strings are splitted at separators such as commas. I am thinking here that a comma is a certain string “,”, and this can be generalized to “splitting” at any substring. With  $\text{\TeX}$ , the analogues are (a) the token with the character code of the comma and category code 12, or the token list consisting of this single token,—and (b) other lists of tokens ...

Anyway, calling a triple  $(\alpha, \sigma, \beta)$  of token lists such that  $\tau = \alpha\sigma\beta$  a “split” of  $\tau$  is not necessarily a bad idea. Moreover, the blank space example (Section 3.1) is very close to the original idea of splitting at separators, a blank space is about as common as a separator as the comma is.

Finally, according to [en.wiktionary.org](http://en.wiktionary.org), the Proto-Indo-European origin of “to bite” just means “to split.”<sup>8</sup> So in  $\text{\TeX}$ 's mouth, splitting and biting is the same.

<sup>6</sup>[tug.org/TUGboat/tb11-2/tb28jeffrey.pdf](http://tug.org/TUGboat/tb11-2/tb28jeffrey.pdf)

<sup>7</sup>[en.wikipedia.org/wiki/String\\_functions#split](http://en.wikipedia.org/wiki/String_functions#split)

<sup>8</sup>[en.wiktionary.org/wiki/bite#Etymology](http://en.wiktionary.org/wiki/bite#Etymology)